

Sibling-First Recursive Graph Drawing for Java Bytecode Md. Jawaherul Alam, Michael T. Goodrich, and Timothy Johnson **Donald Bren School of Information and Computer Sciences** University of California, Irvine

Overview

We describe a tool, the JVM abstracting abstract machine (Jaam) Visualizer, or "J-Viz" for short, which is intended for use by security analysts to find algorithmic vulnerabilities through the exploration of graphs derived from Java bytecode.

Our tool accomplishes this by constructing a graph representing the control flow of the program and displaying it interactively using a canonical ordering we call the *sibling-first recursive ordering*.

We use the following algorithm:

- . Construct the control flow graph using 1-CFA.
- 2. Draw a spanning tree using the sibling-first recursive ordering, putting each branch in a separate lane.
- 3. Draw edges that are not in spanning tree.
- 4. Highlight suspicious areas of the graph.
- 5. Group nodes hierarchically.

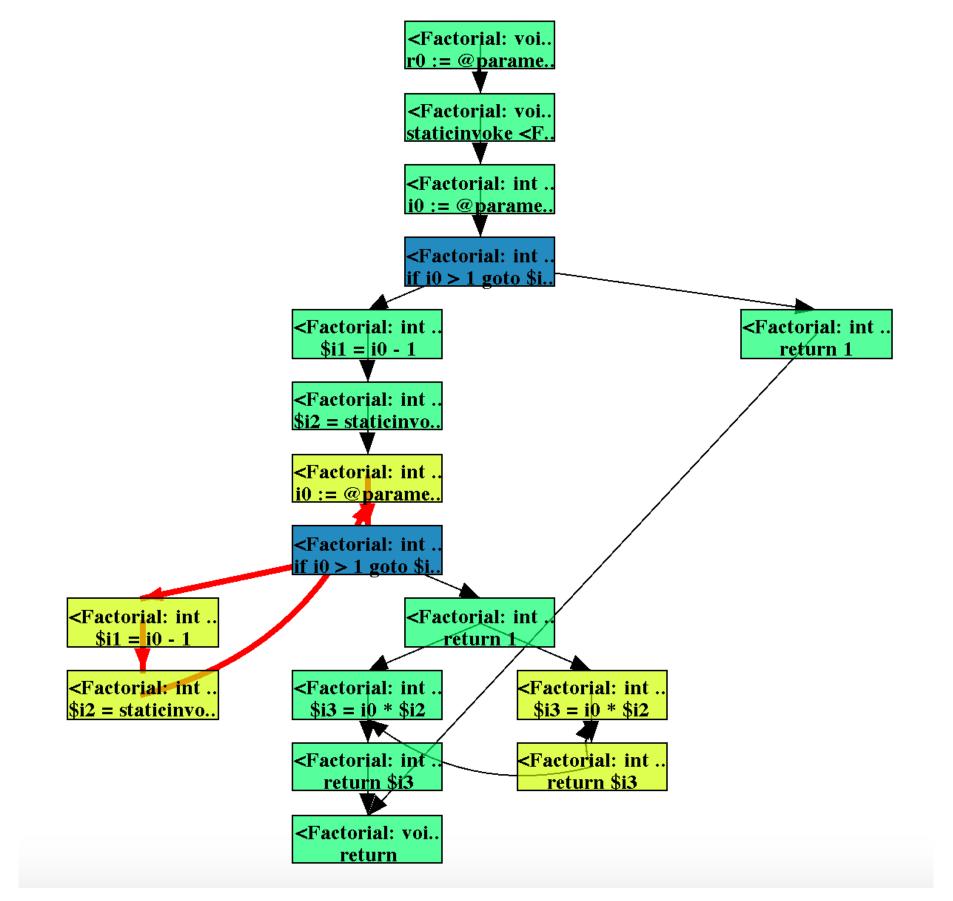
Control Flow Analysis (CFA)

Our desired properties for a control flow graph are:

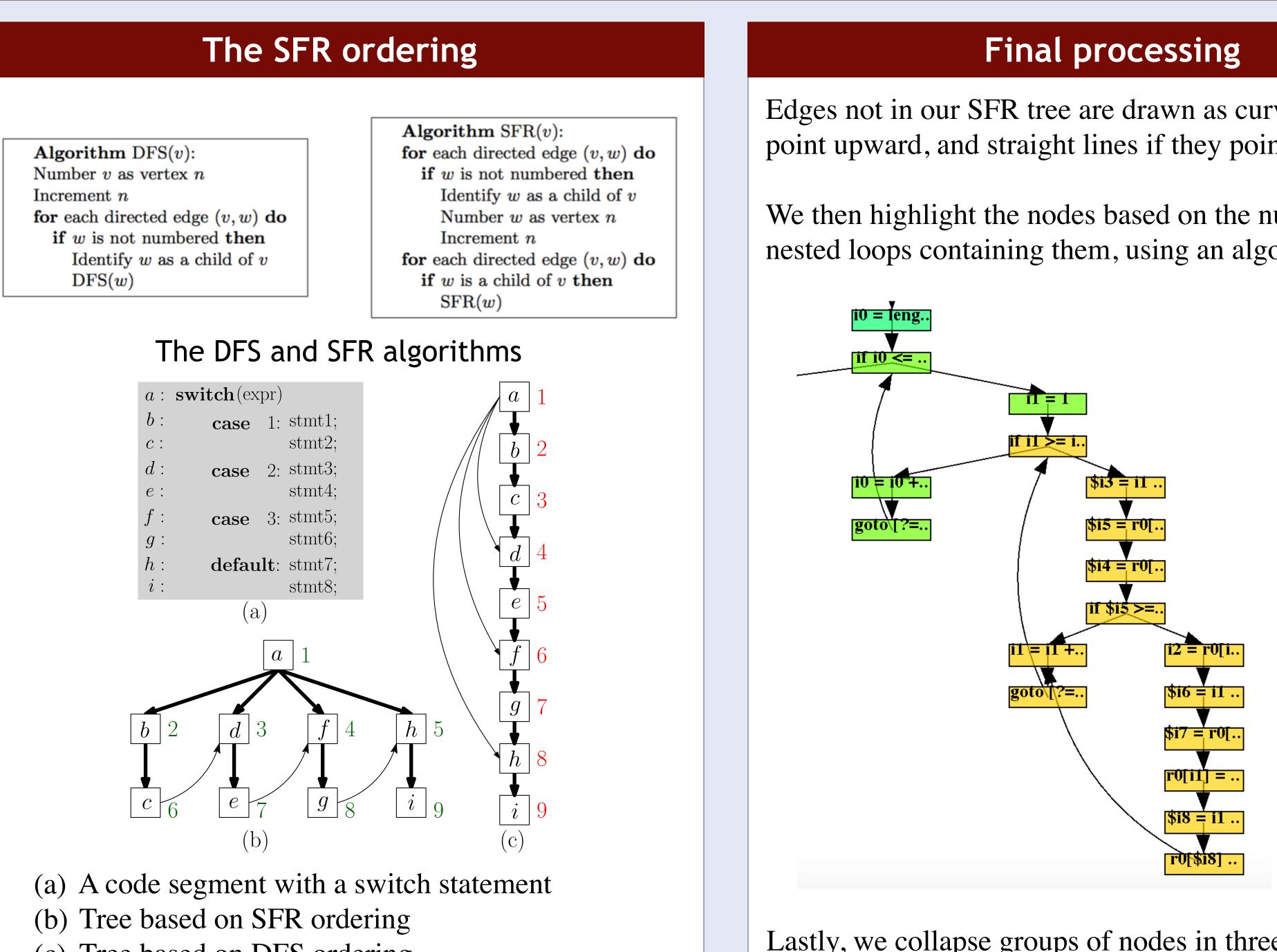
- Soundness: Including every possible execution path
- Precision: Excluding impossible execution paths

The *k*-CFA hierarchy [1] provides a range of algorithms for constructing a sound control flow graph of a program, trading off increased precision against increased size of the graph produced.

For a given k, each state stores a single instruction and the most recent k function calls.



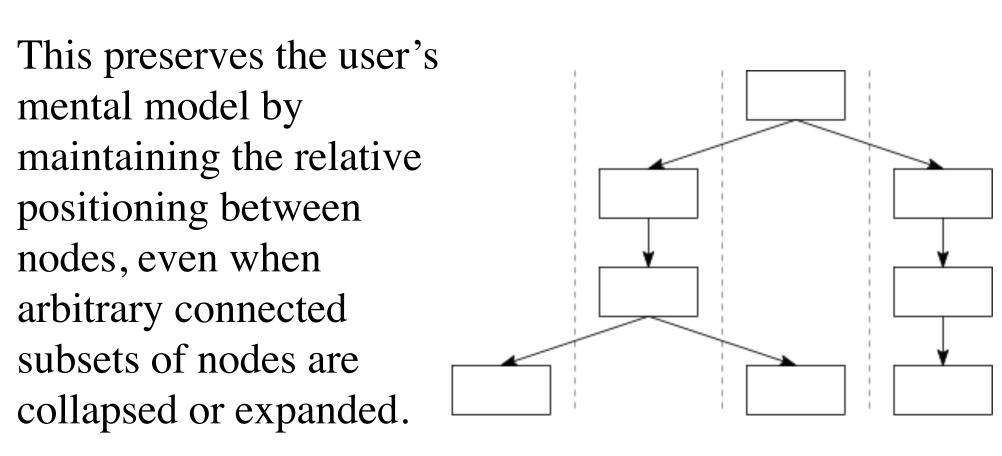
A recursive factorial function analyzed using 1-CFA. The two highlighted nodes represent the same line of code, but are called in different contexts.



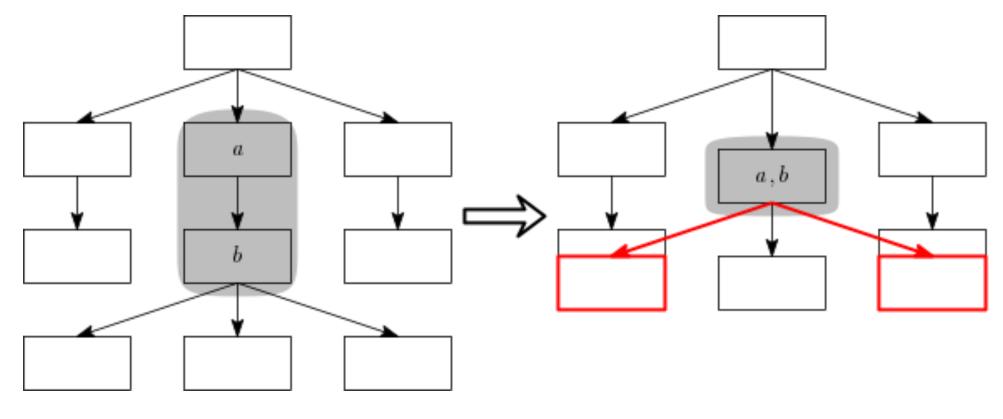
(c) Tree based on DFS ordering

Tree layout

We require that a node can only be drawn directly below another node if it is a descendant in the SFR tree. This divides our graph into lanes, as in the figure below.



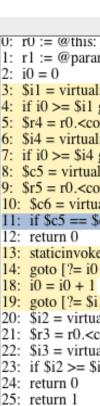
Neglecting this condition can produce collisions, such as in the following example where nodes *a* and *b* are collapsed to a single node.



We then highlight the nodes based on the number of nested loops containing them, using an algorithm from [2].

This program checks a password, one character at a time. But it exits as soon as it finds the first incorrect character, allowing for a timing attack that can determine the password with only a handful of guesses.

This can be found quickly because our graph shows all of the possible exits from the loop.



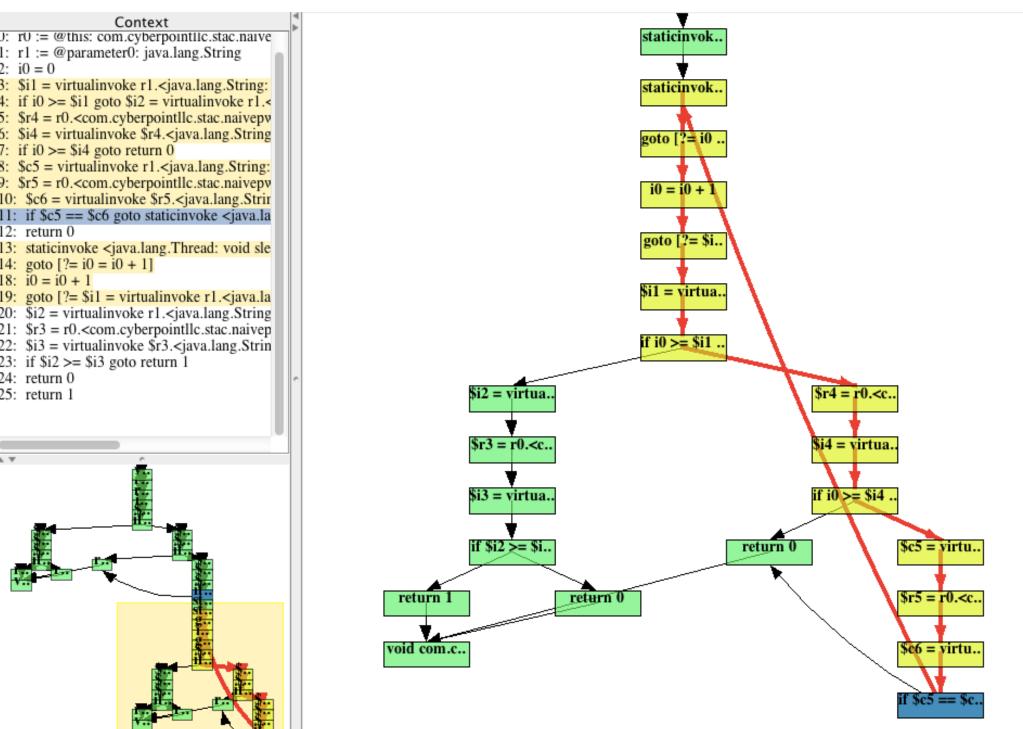


Edges not in our SFR tree are drawn as curves if they point upward, and straight lines if they point downward.

Lastly, we collapse groups of nodes in three ways: - Chains

- Methods
- Chains of methods

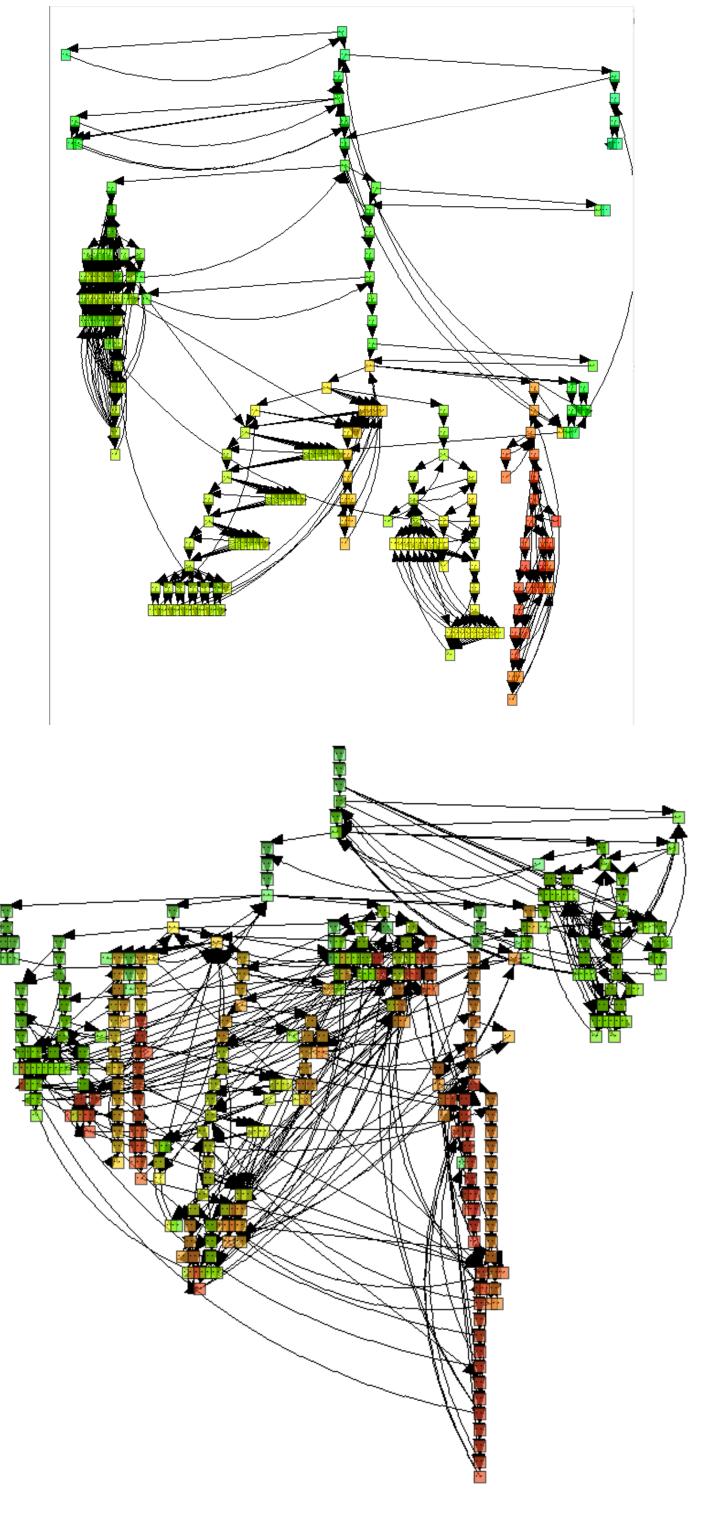
Case study 1



But in each one, the vulnerability appears in the red portion of the graph, allowing the human analyst to check only a small fraction of the code.



The next two programs were provided as part of the DARPA challenge. They are much larger, containing thousands of lines of code.



References

. Shivers, O.G.: Control-Flow Analysis of Higher-Order Languages. Ph.D. thesis, Carnegie Mellon University (1991)

2. Wei, T., Mao, J., Zou, W., Chen, Y.: A new algorithm for identifying loops in decompilation. In: Static Analysis, pp. 170–183. Springer (2007)

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