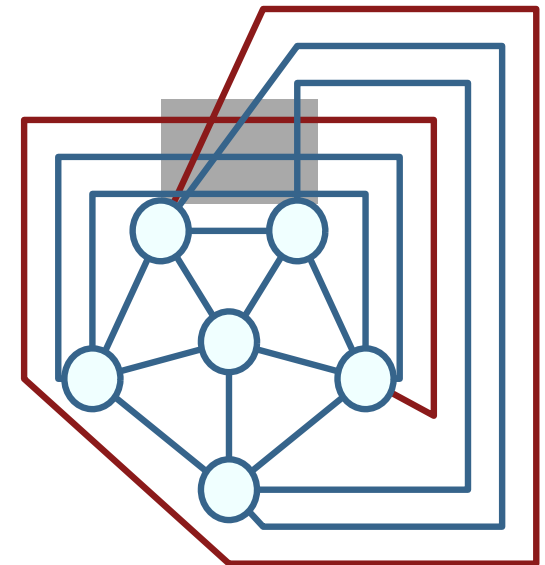
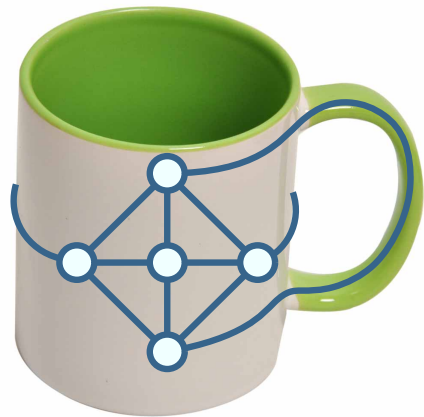


The Bundled Crossing Number

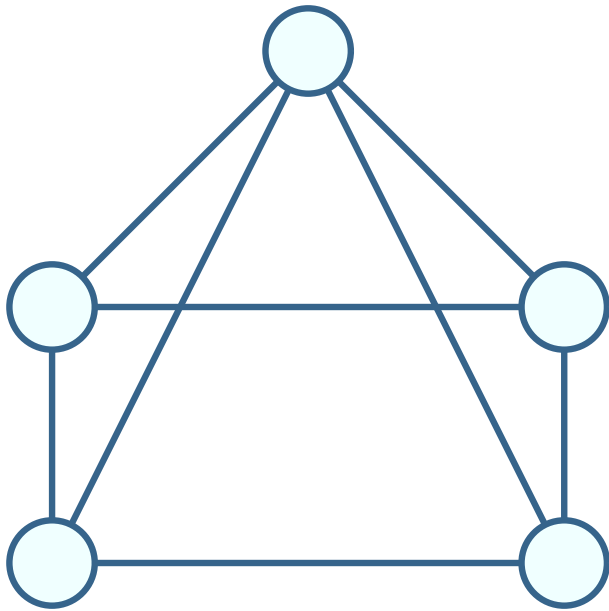
Md. Jawaherul Alam
Martin Fink
Sergey Pupyrev



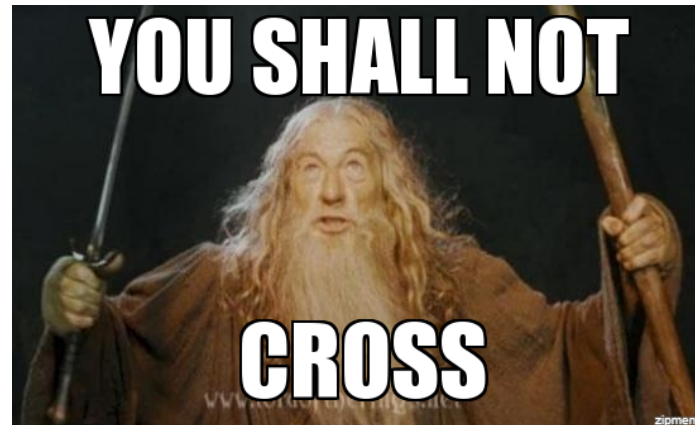
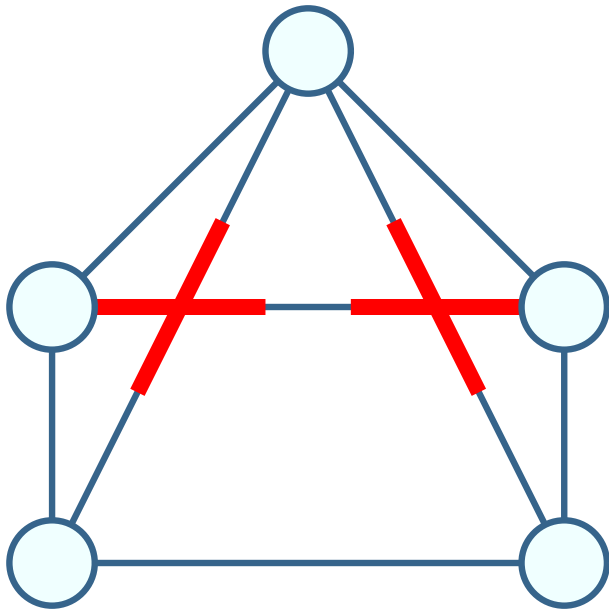
GD 2016 – Athens
September 20



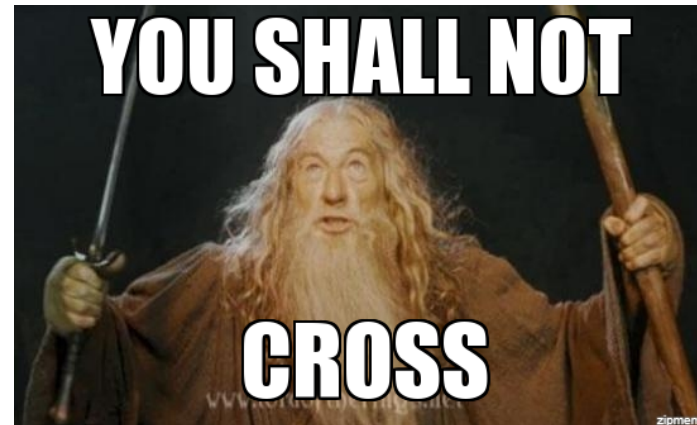
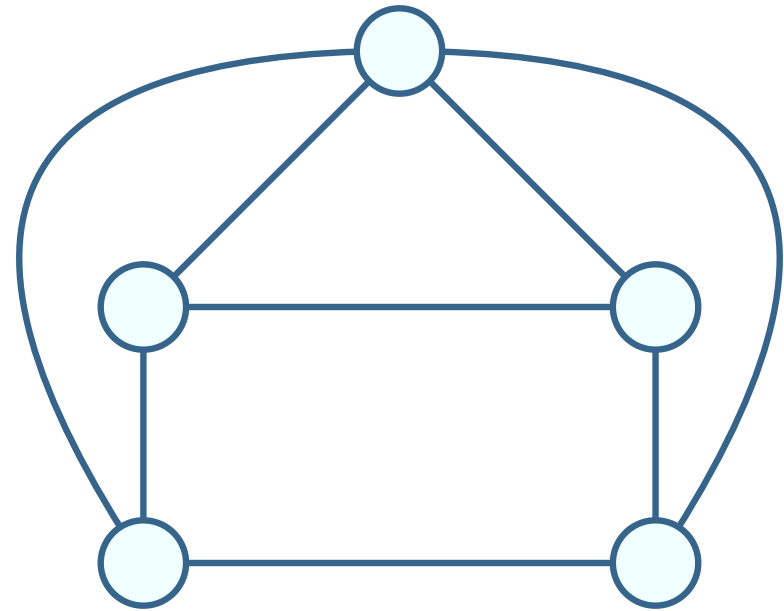
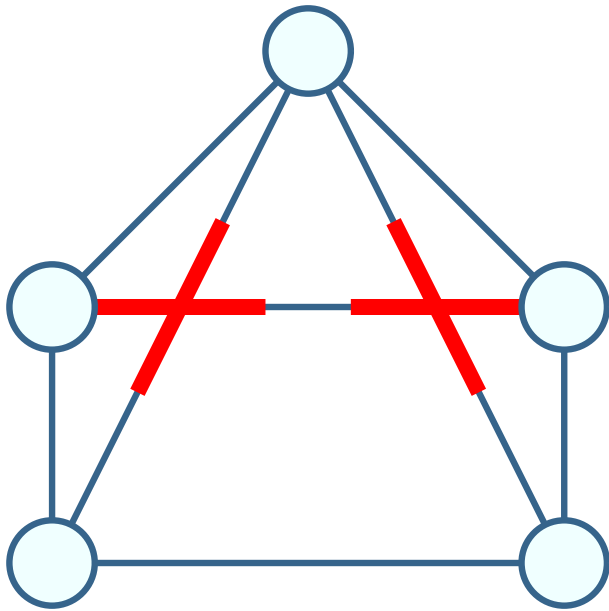
Edge Crossings



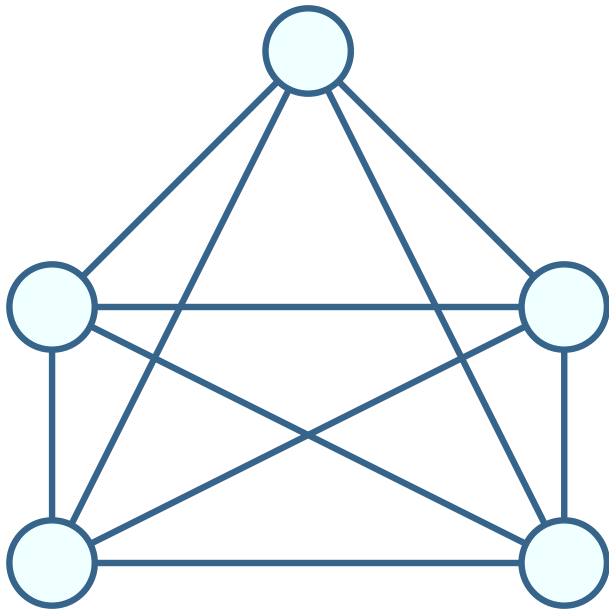
Edge Crossings



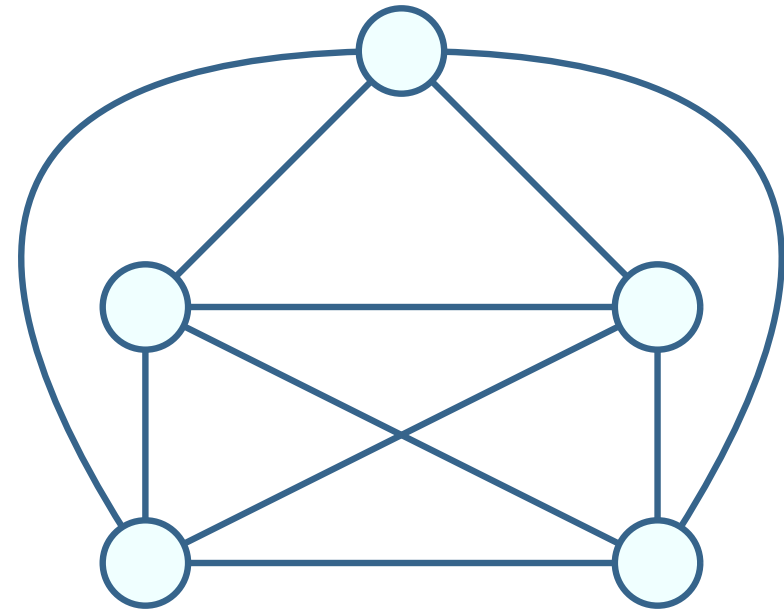
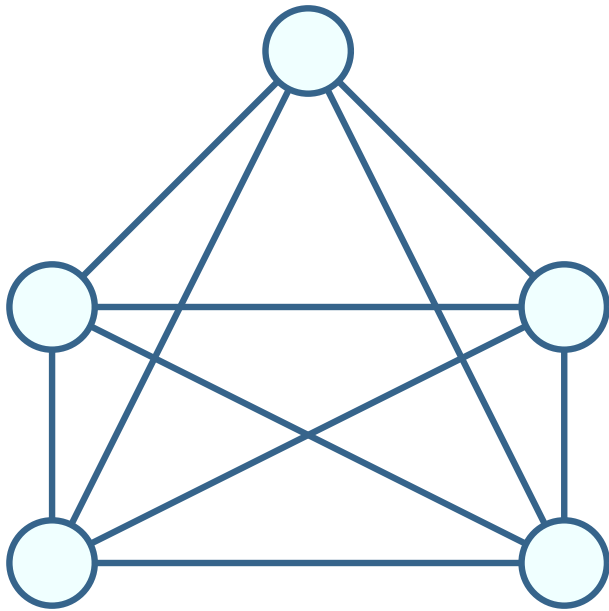
Edge Crossings



Edge Crossings



Edge Crossings



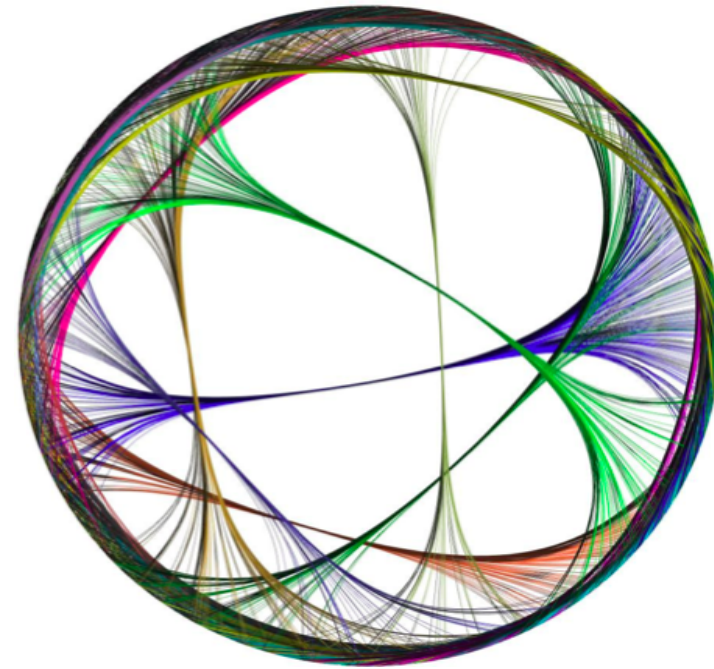
Edge Crossings



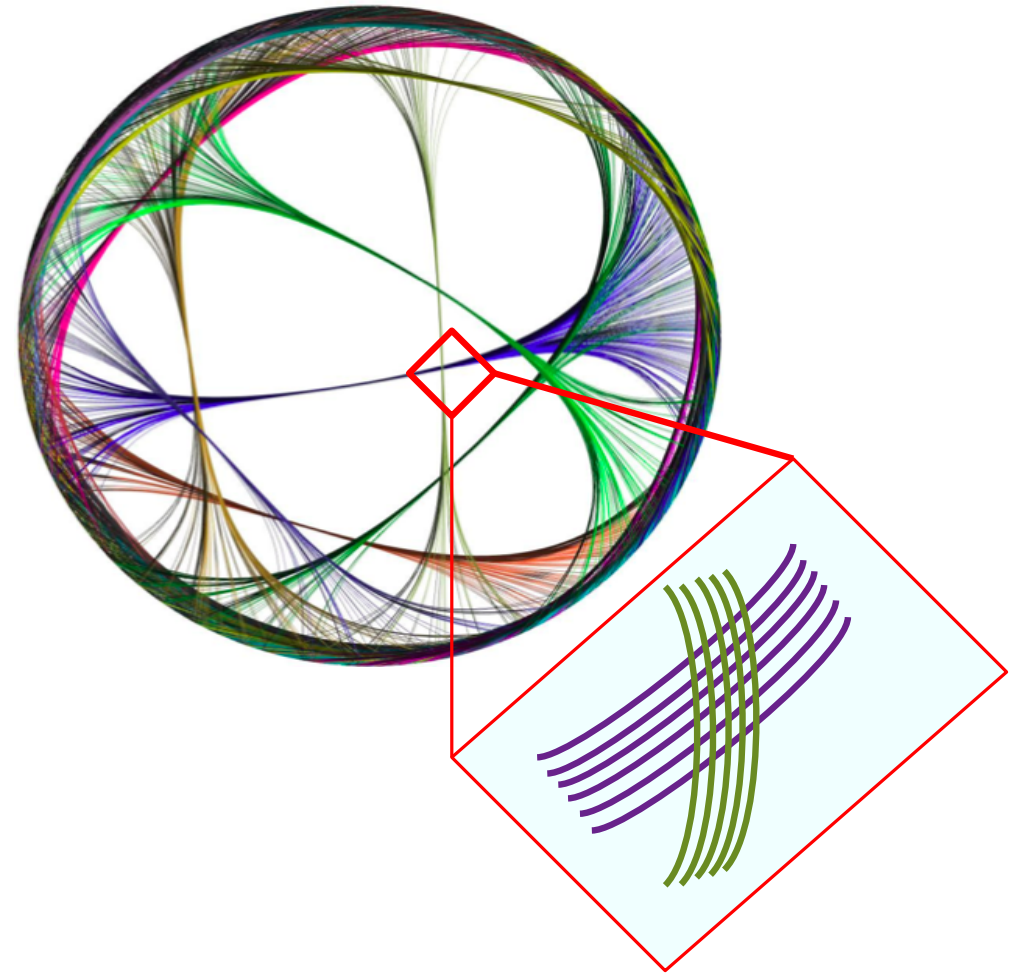
Edge Crossings



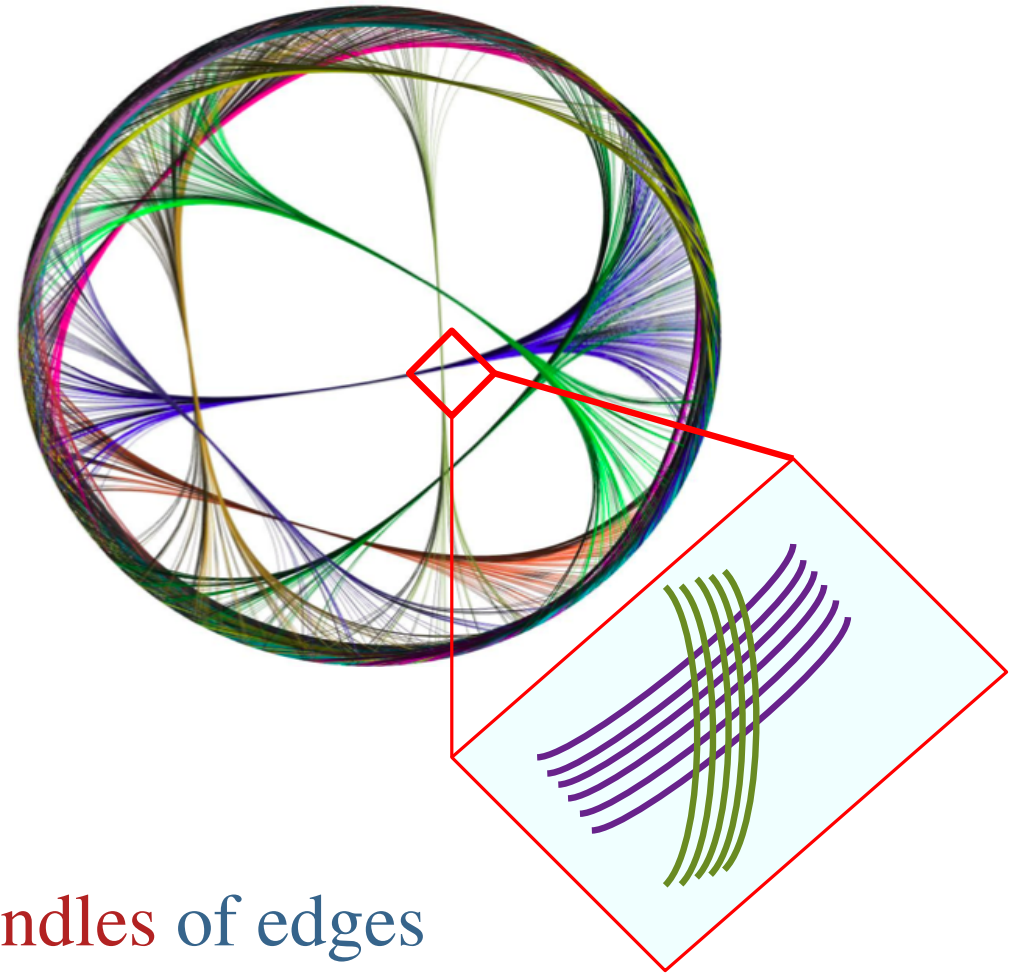
Edge Crossings



Edge Crossings



Bundled Crossings



- Crossings between two **bundles** of edges
- Contained within **disjoint** pseudodisks

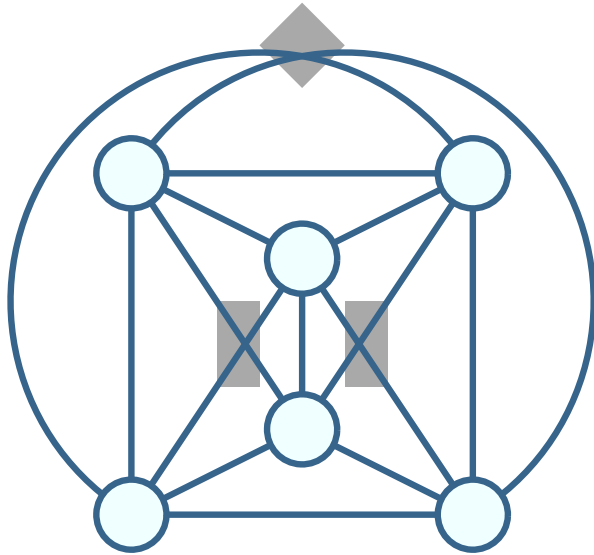


Bundled Crossing Number

minimum number of bundles to group all the crossings



Bundled Crossing Number

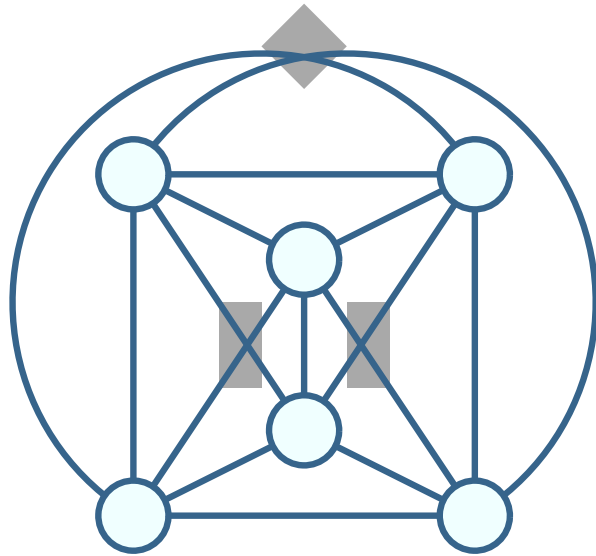


3 bundles

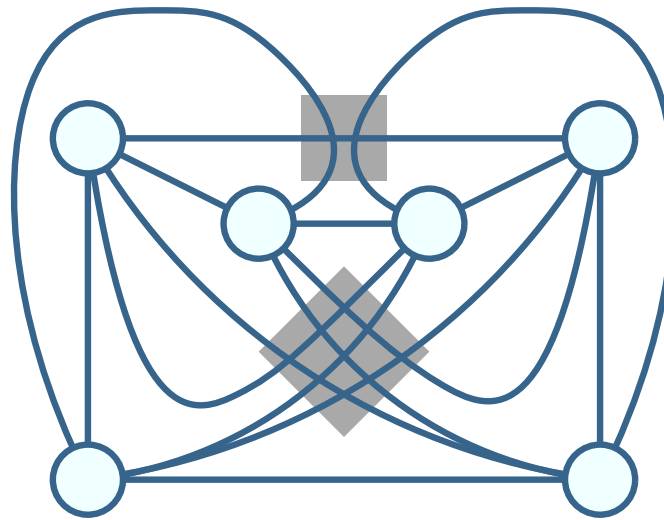
minimum number of bundles to group all the crossings



Bundled Crossing Number



3 bundles

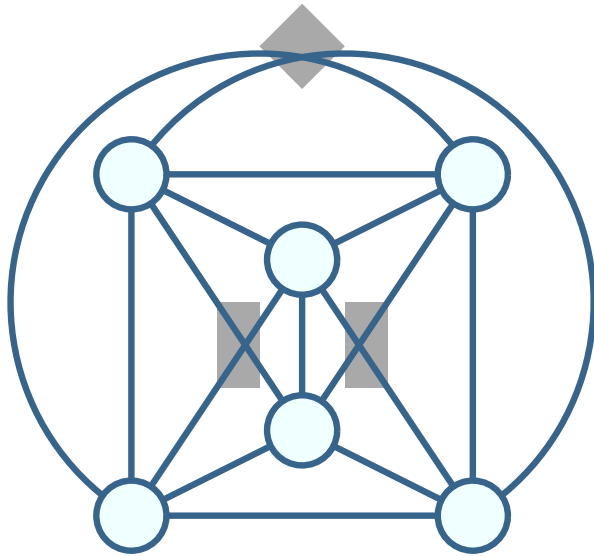


2 bundles

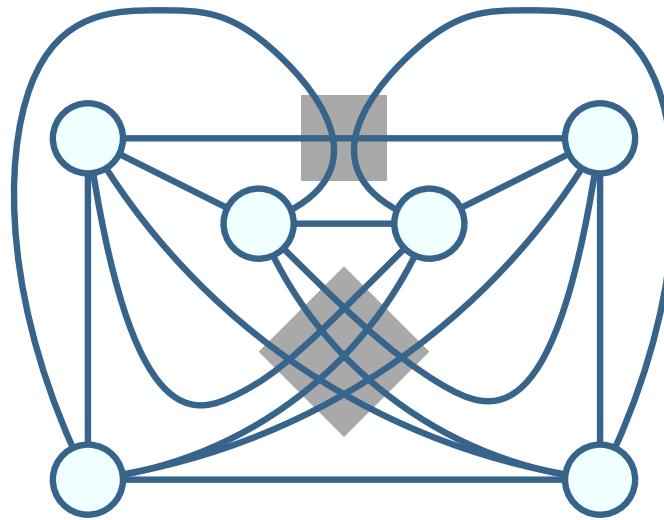
minimum number of bundles to group all the crossings



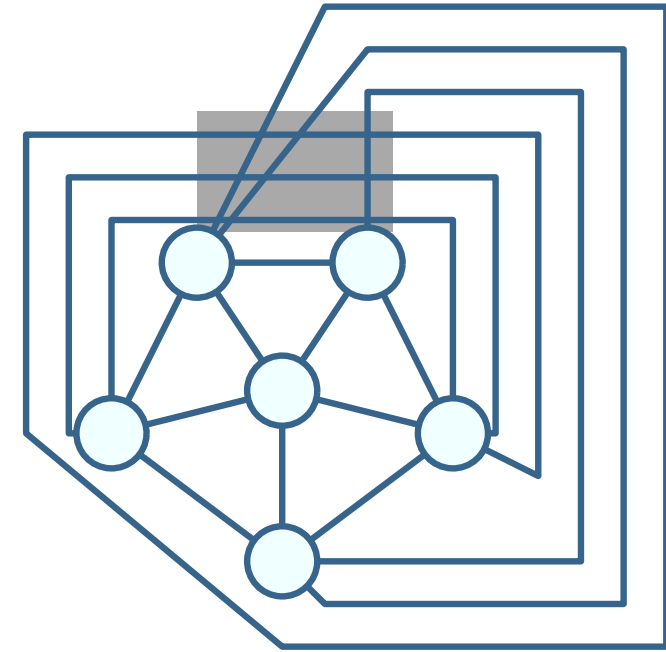
Bundled Crossing Number



3 bundles



2 bundles

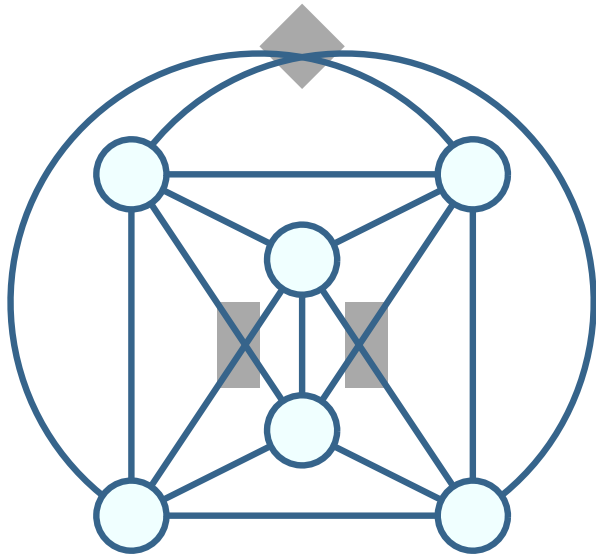


1 bundle

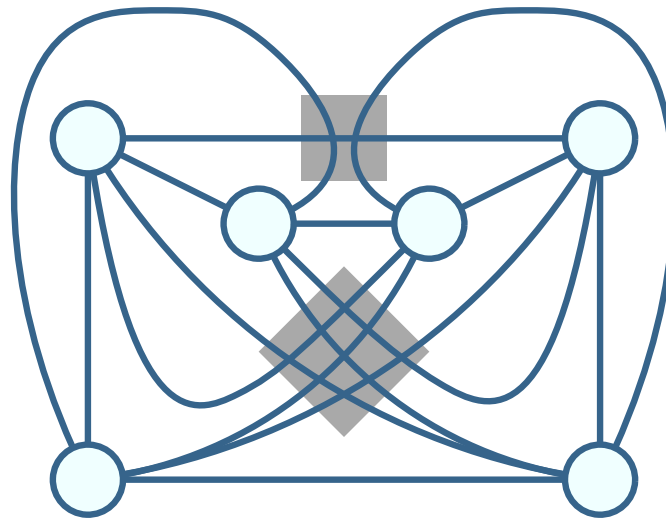
minimum number of bundles to group all the crossings



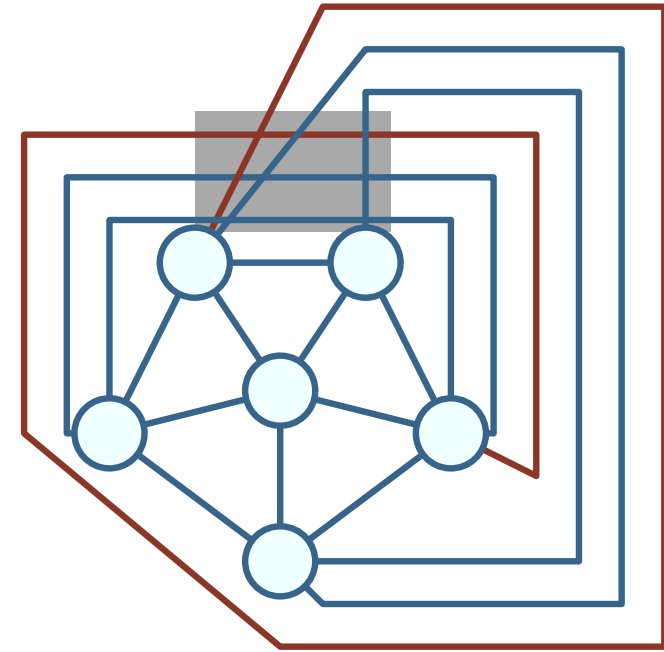
Bundled Crossing Number



3 bundles



2 bundles



1 bundle

minimum number of bundles to group all the crossings



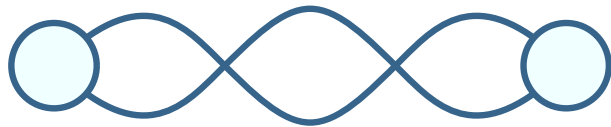
Variants

Unrestricted Drawing

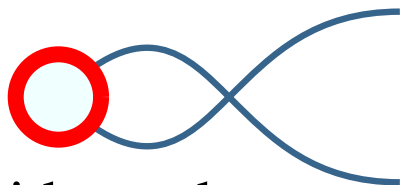
Simple Drawing



self crossing



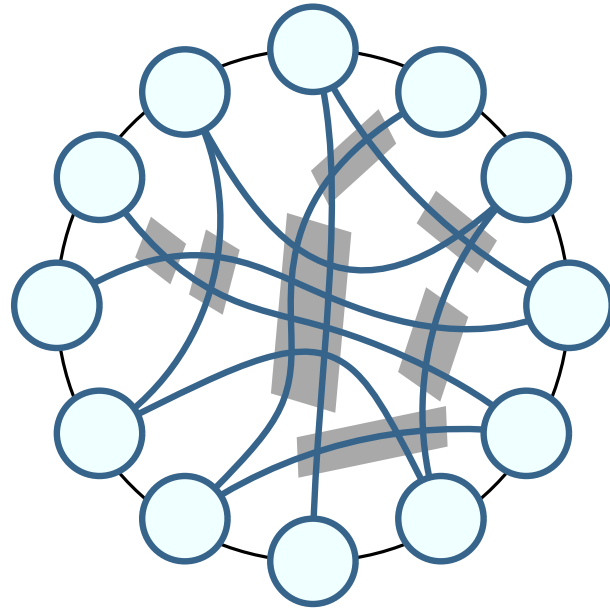
multiple crossings



incident edge crossing



Variants

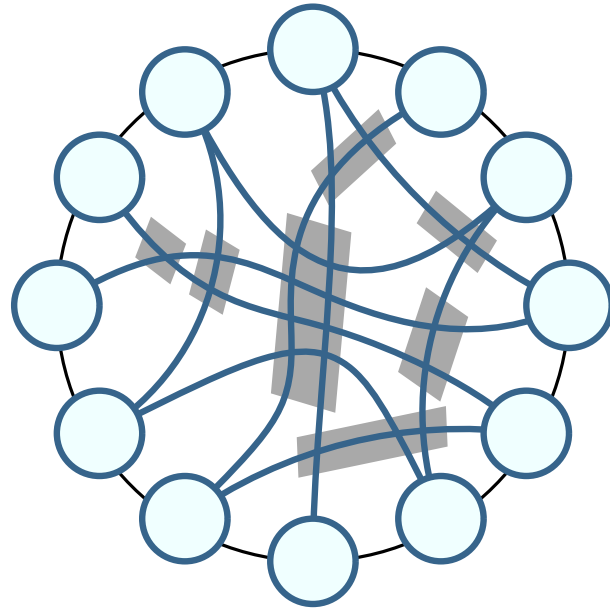


Circular Drawing

- simple drawing
- all vertices on a circle
- all edges inside the circle



Variants



Circular Drawing

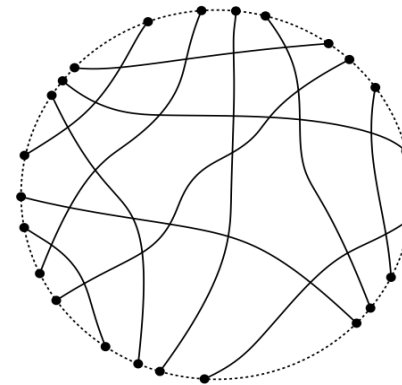
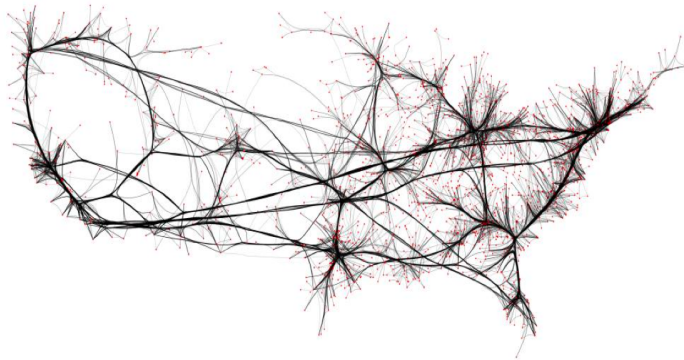
- simple drawing
- all vertices on a circle
- all edges inside the circle
- **Fixed order vs flexible order of vertices**



Related Work

Minimizing Crossings

- **NP-hard**, even for cubic graphs [Cabello 2013]
- No constant approximation unless **$P \neq NP$**
- $O(n^{9/10})$ -approximation for **bounded degree** [Chuzhoy 2011]



Bundled Crossings

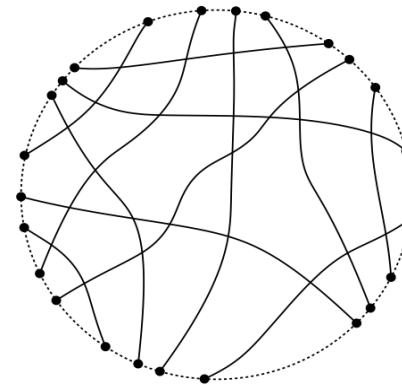
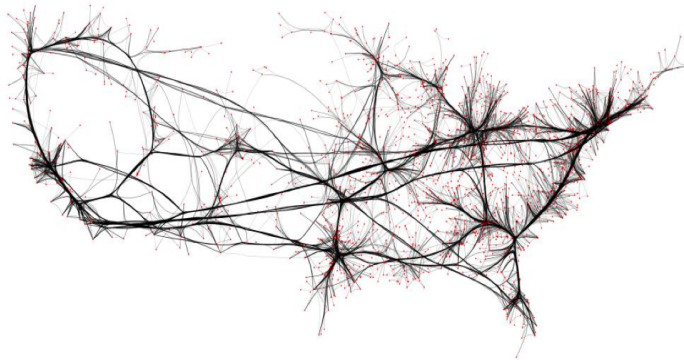
- **NP-hard**, for a **fixed embedding** [Fink et al. 2013]
- constant-factor approximation for **circular layout**



Related Work

Minimizing Crossings

- **NP-hard**, even for cubic graphs [Cabello 2013]
- No constant approximation unless **$P \neq NP$**
- $O(n^{9/10})$ -approximation for **bounded degree** [Chuzhoy 2011]



Bundled Crossings

- **NP-hard**, for a **fixed embedding** [Fink et al. 2013]
- constant-factor approximation for **circular layout**



Our Result

- We address the problem in the **fixed embedding** setting

	Lower Bound	Upper Bound	Approximation
Unrestricted	g	g	--
Simple	g	--	--
	$\frac{(m - (3n - 6))}{6}$	$m - 1$	$\frac{6c}{(c-3)}, m > cn$
circular (flexible)	$\frac{(m - (2n - 3))}{6}$	$m - 1$	$\frac{6c}{(c-2)}, m > cn$
circular (fixed)	$m/16$	$m - 1$	16



Our Result

- We address the problem in the **fixed embedding** setting

	Lower Bound	Upper Bound	Approximation
Unrestricted	g	g	--
Simple	g	--	--
	$\frac{(m-(3n-6))}{6}$	$m - 1$	$\frac{6c}{(c-3)}, m > cn$
circular (flexible)	$\frac{(m-(2n-3))}{6}$	$m - 1$	$\frac{6c}{(c-2)}, m > cn$
circular (fixed)	$m/16$	$m - 1$	16



Our Result

- We address the problem in the **fixed embedding** setting

	Lower Bound	Upper Bound	Approximation
Unrestricted	g	g	--
Simple	g	--	--
	$\frac{(m - (3n - 6))}{6}$	$m - 1$	$\frac{6c}{(c-3)}, m > cn$
circular (flexible)	$\frac{(m - (2n - 3))}{6}$	$m - 1$	$\frac{6c}{(c-2)}, m > cn$
circular (fixed)	$m/16$	$m - 1$	16



Our Result

- We address the problem in the **fixed embedding** setting

	Lower Bound	Upper Bound	Approximation
Unrestricted	g	g	--
Simple	g	--	--
	$\frac{(m - (3n - 6))}{6}$	$m - 1$	$\frac{6c}{(c-3)}, m > cn$
circular (flexible)	$\frac{(m - (2n - 3))}{6}$	$m - 1$	$\frac{6c}{(c-2)}, m > cn$
circular (fixed)	$m/16$	$m - 1$	16



Bundled Crossing Number and Genus

Theorem

For unrestricted drawings of G , $bc(G) = g(G)$



Bundled Crossing Number and Genus

Theorem

For unrestricted drawings of G , $bc(G) = g(G)$

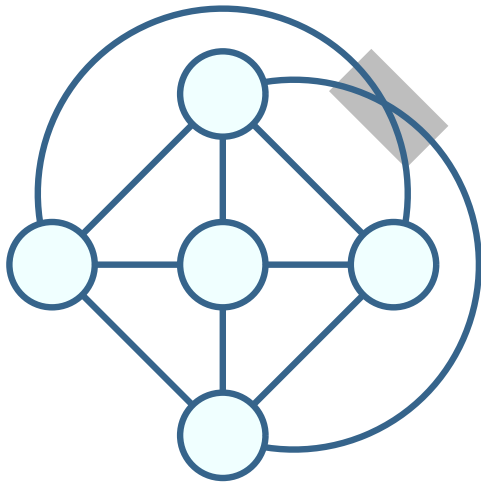


Bundled Crossing Number and Genus

Theorem

For unrestricted drawings of G , $bc(G) = g(G)$

Proof [$bc(G) \geq g(G)$]:

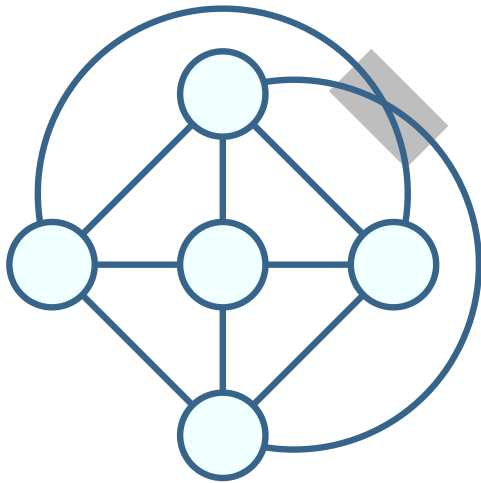


Bundled Crossing Number and Genus

Theorem

For unrestricted drawings of G , $bc(G) = g(G)$

Proof [$bc(G) \geq g(G)$]:



For each bundle, create a handle and re-route

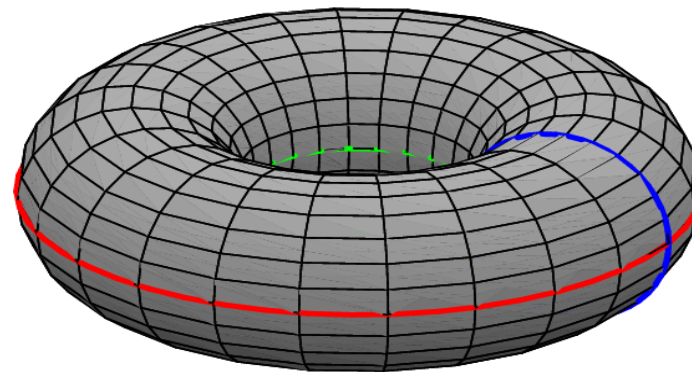
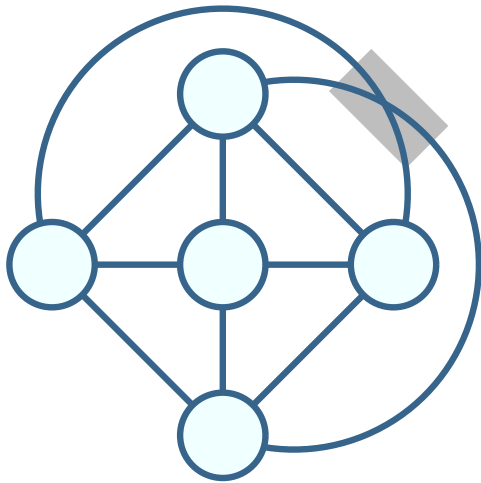


Bundled Crossing Number and Genus

Theorem

For unrestricted drawings of G , $bc(G) = g(G)$

Proof [$bc(G) \geq g(G)$]:



For each bundle, create a handle and re-route

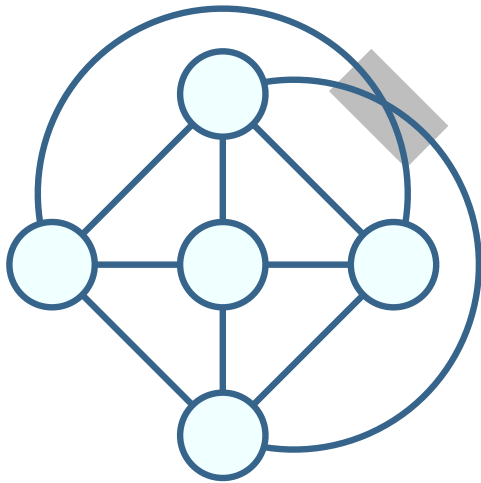


Bundled Crossing Number and Genus

Theorem

For unrestricted drawings of G , $bc(G) = g(G)$

Proof [$bc(G) \geq g(G)$]:



For each bundle, create a handle and re-route

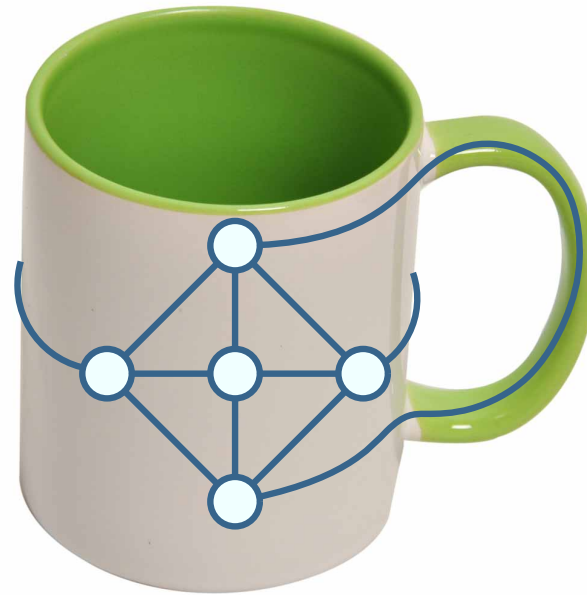
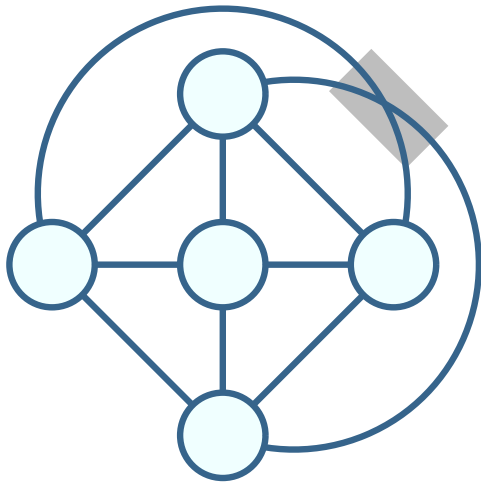


Bundled Crossing Number and Genus

Theorem

For unrestricted drawings of G , $bc(G) = g(G)$

Proof [$bc(G) \geq g(G)$]:



For each bundle, create a handle and re-route

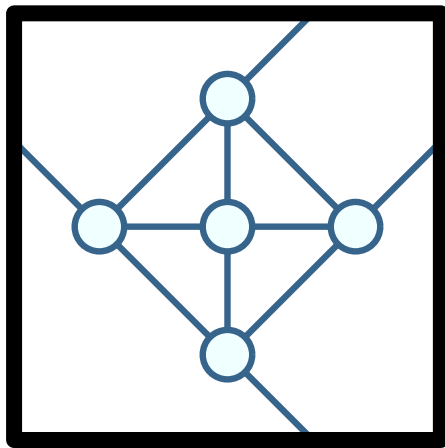


Bundled Crossing Number and Genus

Theorem

For unrestricted drawings of G , $bc(G) = g(G)$

Proof [$bc(G) \leq g(G)$]:

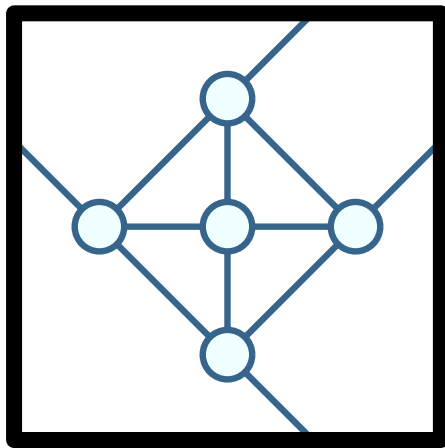


Bundled Crossing Number and Genus

Theorem

For unrestricted drawings of G , $bc(G) = g(G)$

Proof [$bc(G) \leq g(G)$]:



re-route each edge touching the boundary through outside

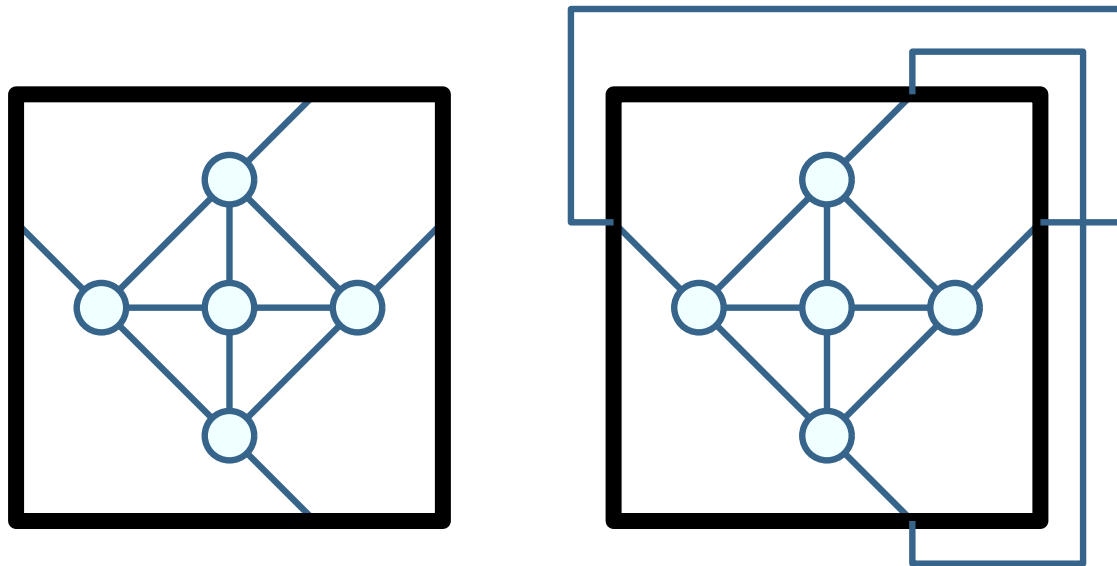


Bundled Crossing Number and Genus

Theorem

For unrestricted drawings of G , $bc(G) = g(G)$

Proof [$bc(G) \leq g(G)$]:



re-route each edge touching the boundary through outside

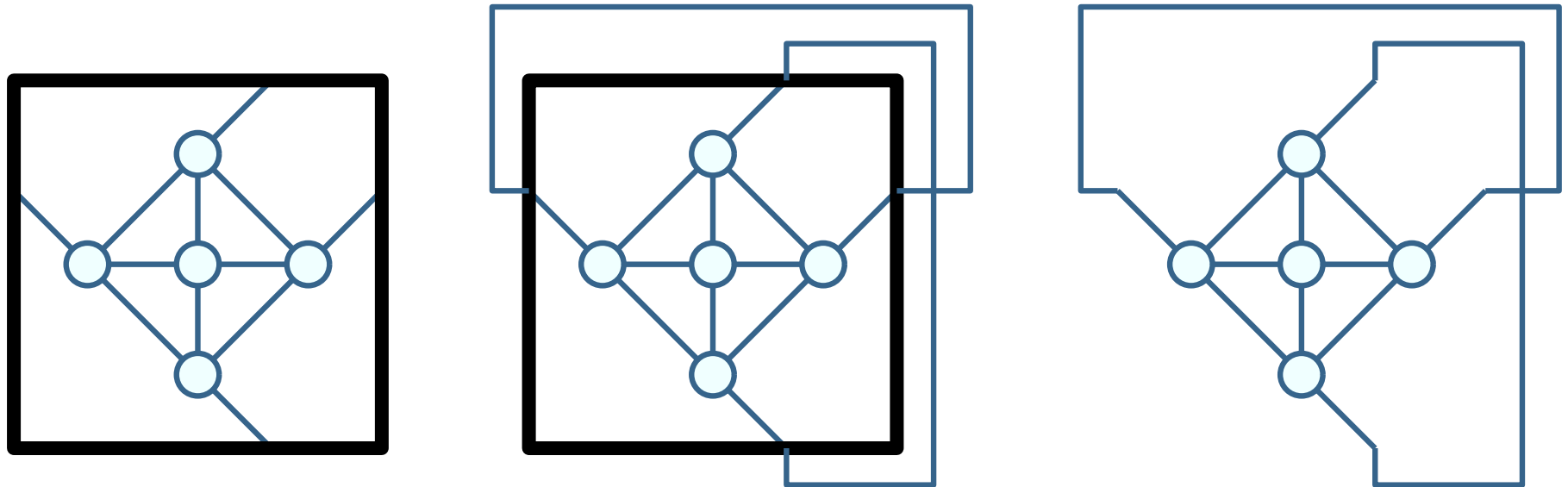


Bundled Crossing Number and Genus

Theorem

For unrestricted drawings of G , $bc(G) = g(G)$

Proof [$bc(G) \leq g(G)$]:



re-route each edge touching the boundary through outside

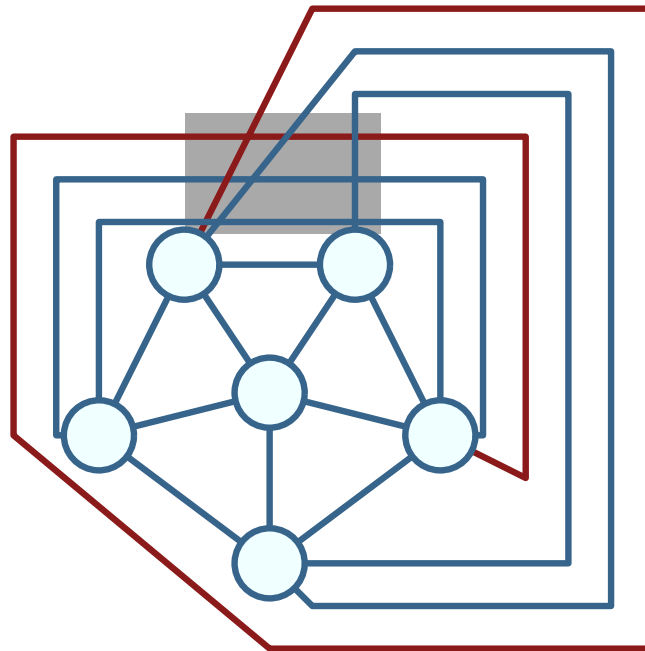


Bundled Crossing Number and Genus

Theorem

For simple drawings of G , $bc(G) \geq g(G)$.

There is some G for which $bc(G) > g(G)$.



Our Result

- We address the problem in the **fixed embedding** setting

	Lower Bound	Upper Bound	Approximation
Unrestricted	g	g	--
Simple	g	--	--
	$\frac{(m - (3n - 6))}{6}$	$m - 1$	$\frac{6c}{(c-3)}, m > cn$
circular (flexible)	$\frac{(m - (2n - 3))}{6}$	$m - 1$	$\frac{6c}{(c-2)}, m > cn$
circular (fixed)	$m/16$	$m - 1$	16



Our Result

- We address the problem in the **fixed embedding** setting

	Lower Bound	Upper Bound	Approximation
Unrestricted	g	g	--
Simple	g	--	--
	$\frac{(m - (3n - 6))}{6}$	$m - 1$	$\frac{6c}{(c-3)}, m > cn$
circular (flexible)	$\frac{(m - (2n - 3))}{6}$	$m - 1$	$\frac{6c}{(c-2)}, m > cn$
circular (fixed)	$m/16$	$m - 1$	16



Bundled Crossing in Circular Layout

Theorem

For fixed order circular layout of G , $bc(G) \geq m/16$.

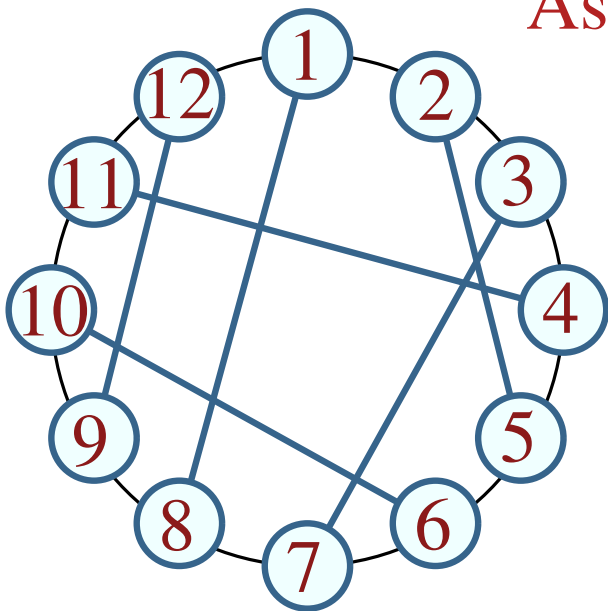


Bundled Crossing in Circular Layout

Theorem

For fixed order circular layout of G , $bc(G) \geq m/16$.

Assume that the edges form a matching

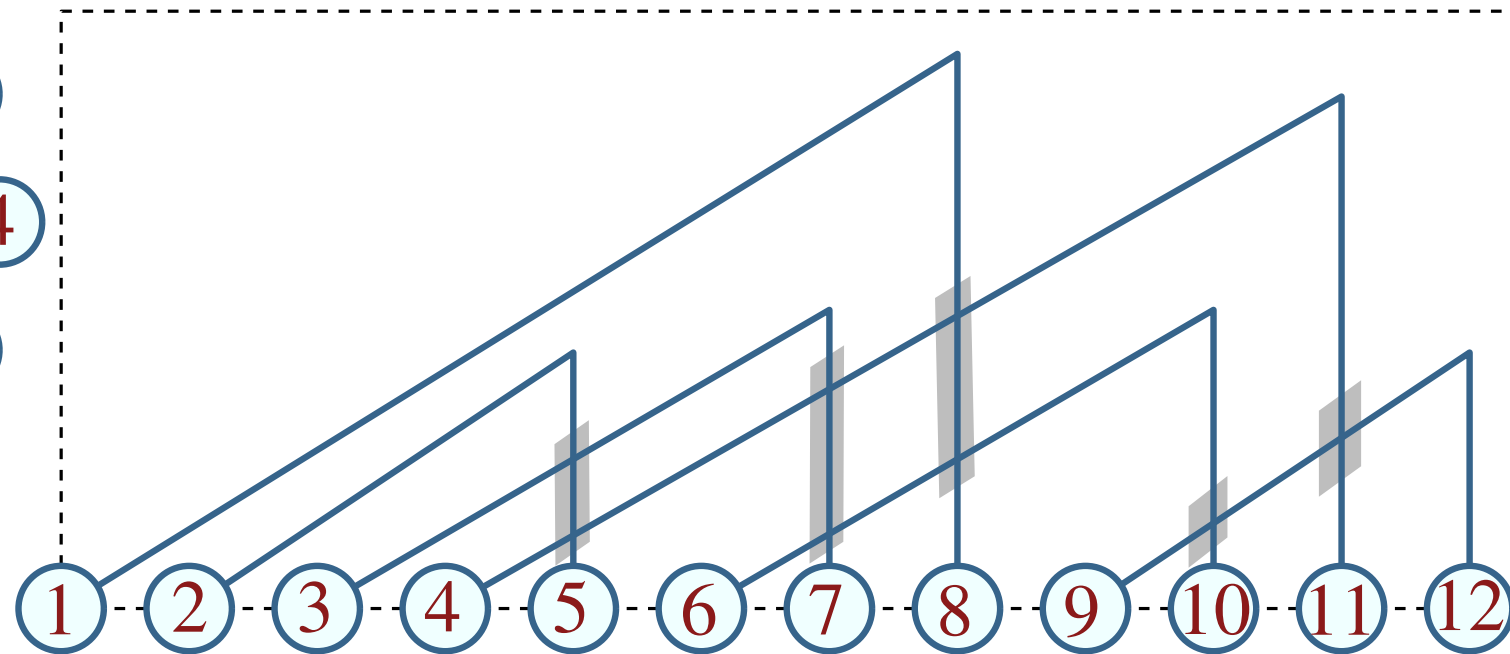
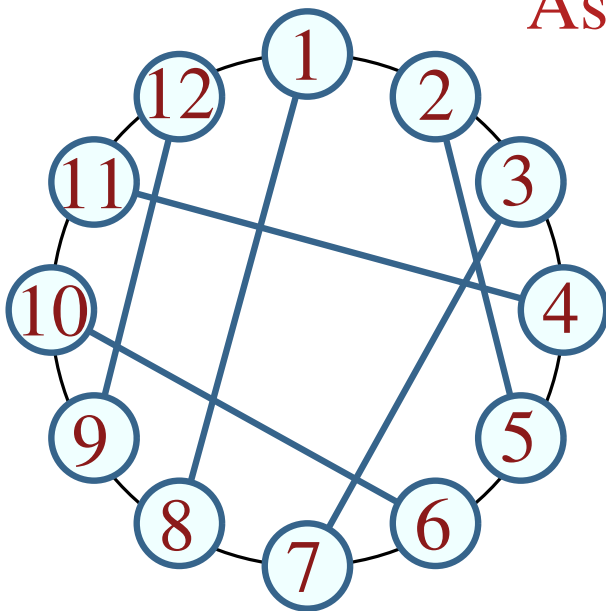


Bundled Crossing in Circular Layout

Theorem

For fixed order circular layout of G , $bc(G) \geq m/16$.

Assume that the edges form a matching



Summary

	Lower Bound	Upper Bound	Approximation
Unrestricted	g	g	--
Simple	g	--	--
	$\frac{(m - (3n - 6))}{6}$	$m - 1$	$\frac{6c}{(c-3)}, m > cn$
circular (flexible)	$\frac{(m - (2n - 3))}{6}$	$m - 1$	$\frac{6c}{(c-2)}, m > cn$
circular (fixed)	$m/16$	$m - 1$	16



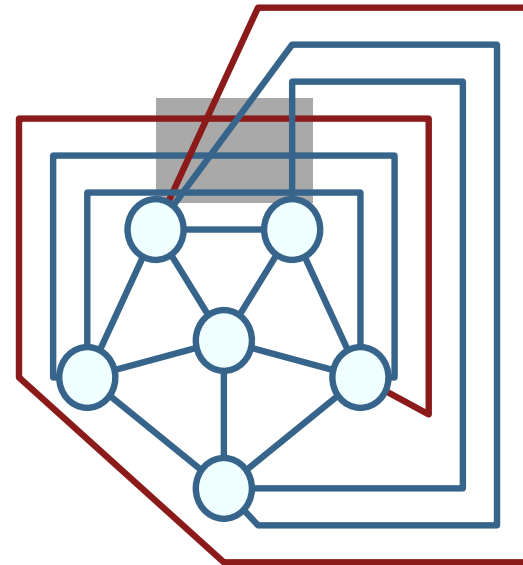
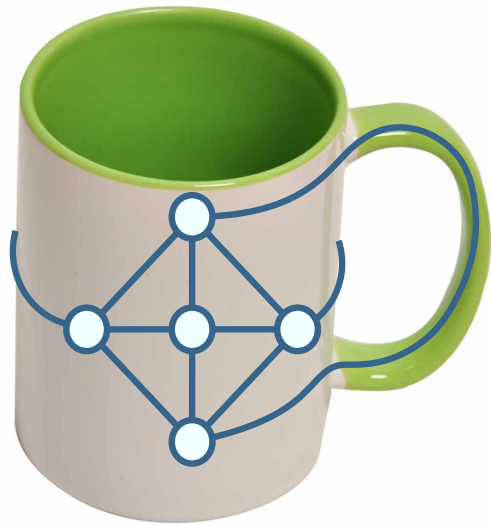
Summary

	Lower Bound	Upper Bound	Approximation
Unrestricted	g	g	--
Simple	g	--	--
	$\frac{(m - (3n - 6))}{6}$	$m - 1$	$\frac{6c}{(c-3)}, m > cn$
circular (flexible)	$\frac{(m - (2n - 3))}{6}$	$m - 1$	$\frac{6c}{(c-2)}, m > cn$
circular (fixed)	$m/16$	$m - 1$	16



Future Work

- Circular bundled crossing: complexity results
 - **NP**-hard?
- Better approximation for sparse graphs



Future Work

- Circular bundled crossing: complexity results
 - **NP-hard?**
- Better approximation for sparse graphs

